Deductive Verification of Realistic OCaml Code*

Carlos Pinto¹, Mário Pereira², and Simão Melo de Sousa¹

 $^1\,$ NOVA LINCS, Universidade da Beira Interior, Portugal $^2\,$ NOVA LINCS, Nova School of Science and Technology, Portugal

Introduction. Cameleer [6] is a deductive verification platform for OCaml programs, actively developed over the last year, built on top of the Why3 framework [4]. It leverages on the recently proposed GOSPEL [1] (Generic OCaml SPEcification Language), which allows programmers to attach rigorous, yet readable, behavioral specification to OCaml code. GOSPEL has been already used in a framework for extracting correctby-construction OCaml implementations with respect to provided, annotated interfaces [2]. With Cameleer, we have pushed this approach further, by centering the verification effort directly on the OCaml implementation. We believe Cameleer has now reached enough maturity to tackle the verification of realistic OCaml code. This presentation proposal reports this claim, namely on our experience of using the Cameleer tool to provide a complete proof of the Set module from the OCaml standard library³.

Deductive verification framework. The workflow of the verification toolchain we report is depicted as follows:



Fig. 1. Cameleer verification workflow.

Cameleer uses the GOSPEL toolchain, in order to parse and manipulate (via the ppxlib library) the abstract syntax tree of GOSPEL-annotated OCaml programs. Cameleer translates the decorated AST into an equivalent WhyML representation, which is then fed to Why3. The user never needs to manipulate the generated WhyML program. In short, the Cameleer user intervenes in the beginning and in the end of the process, *i.e.*, in the initial specifying phase and in the last step, helping Why3 to close the proof.

Module Set. In order to assess the maturity of the Cameleer tool, several OCaml modules implementing data structures and algorithms have been verified. We focus on this presentation on the verification of the Set module of the standard library. This is an interesting module, since it is a realistic piece of OCaml code, widely-used, and presenting subtle design choices.

The following, is an excerpt of the functorial implementation of Set:

```
module type OrderedType = sig
  type t
  val[@logic] compare : t -> t -> int
  (*@ axiom is_pre_order: is_pre_order compare *)
end
```

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³ https://ocaml.org/api/Set.html

```
module Make(Ord: OrderedType) = struct
  let bal l v r =
    . . .
  (*0 res = bal l v r
        requires avl 1 && avl r
        requires abs (height 1 - height r) <= 3</pre>
        requires forall y. mem y 1 -> Ord.compare y v < 0</pre>
        requires forall x. mem x r -> Ord.compare v x < 0
        requires forall x y. mem x 1 -> mem y r -> Ord.compare x y < 0
        ensures avl res
        ensures forall w. w <> v -> occ w res = occ w l + occ w r
                 occ v res = 1
        ensures
                      max (height 1) (height r) <= height res <=</pre>
        ensures
                  1 + max (height 1) (height r)
        ensures abs (height 1 - height r) <= 2 -> res = create 1 v r *)
```

The rebalancing function **bal** is here presented together with a suitable GOSPEL contract. Cameleer and GOSPEL allows the programmer to define logical definitions, for instance to describe what is a well-balanced binary search tree, according to the balance criteria used in the Set module.

Further details will be provided during the presentation itself. In a few words, the **bal** function takes as arguments two balanced trees, 1 and r, a new element v and builds a new balanced tree containing v and all the elements from 1 and r. For such a function it is required for 1 and r to be indeed balanced trees (this is captured by the recursive predicate av1). We use the occ recursive logical function to count the number of occurrences of a certain element inside a tree. Such a predicate allows us to write very simple and elegant contracts, such as all the elements distinct from v maintain their occurrence in the resulting tree, while the occurrence of v is exactly one. The other specification elements have the usual semantics. It is worth mentioning that the informal comments provided in the set.ml source file are very clear and can be, almost directly, translated into GOSPEL specification.

Conclusion. Our proof of a subset of the operations from the Set module is publicly available⁴. It amounts to 122 lines of code and 117 lines of GOSPEL specification. Cameleer automatically feeds to Why3 a translation of the selected Set functions, which generates a total of 521 Verification Conditions (VCs). The entire proof can be replayed in an average time of 66 seconds, measured on a laptop equipped with macOS Catalina, with 8 GB of Ram and 2 CPU cores (2.6 GHz).

The proof effort was conducted mainly by an undergraduate student, with no previous experience on formal verification. This was intentional, in order to assess our claim on the maturity and ease of use of the Cameleer toolchain. We honestly believe our experience is an important contribution towards the effort of building a corpus of realistic, formally verified OCaml modules. Regarding previous verification efforts of AVL data structures, in particular the work by Filliâtre and Letouzey [3], our work clearly improves on proof automation. We believe this is an important aspect, since automated verification approaches are more likely to be adopted by regular programmers, who are not necessarily proof experts. Moreover, the Cameleer approaches precludes the need of OCaml code extraction (which is the approach followed by Coq and Why3, for instance), and for programmers to write entire code bases for the sake of verification.

As future work, we intend to specify and prove higher-order functions from module Set. As a first step, we will tackle the verification of higher-order computation where is reasonable to assume the argument function is effect-free (*e.g.*, map and fold). Finally, a more challenging approach is the proof of higher-order functions with effects, for instance iter. Currently, neither Why3 or GOSPEL support effectful higher-order computation. We plan to follow the approach by Filliâtre and Pereira [5] for the modular treatment of iteration structures, and extend the core Cameleer to cope with the proposed translation of iter-like routines.

⁴ https://github.com/ocaml-gospel/cameleer/blob/testing_carlos/examples/Set.ml

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